

Full Length Research Paper

# Some structural compatibilities of pre $A^*$ -algebra

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**A pre  $A^*$ -algebra is the algebraic form of the 3-valued logic. In this paper, we define a binary operation  $\oplus$  on pre  $A^*$ -algebra and show that  $\langle A, \oplus \rangle$  is a semilattice. We also prove some results on the partial ordering  $\leq_{\oplus}$  which is induced from the semilattice  $\langle A, \oplus \rangle$ . We derive necessary and sufficient conditions for pre  $A^*$ -algebra  $(A, \wedge, \vee, (-)')$  to become a Boolean algebra in terms of this partial ordering and binary operation and also find the necessary conditions for  $(A, \leq_{\oplus})$  as a lattice.**

**Key words:** Pre  $A^*$ - algebra, poset, semilattice, centre, Boolean algebra.

## INTRODUCTION

In 1948 the study of lattice theory was first introduced by Birkhoff (1948). In a draft paper by Manes (1989), the Equational Theory of Disjoint Alternatives, E. G. Maines introduced the concept of ADA (Algebra of disjoint alternatives)  $(A, \wedge, \vee, (-)', (-)_{\pi}, 0, 1, 2)$  which however differs from the definition of the ADA of his later paper (Manes, 1993); ADAs and the equational theory of if-then-else in 1993. While the ADA of the earlier draft seems to be based on extending the if-then-else concept more on the basis of Boolean algebra and the later concept was based on C-algebra  $(A, \wedge, \vee, ')$  introduced by Fernando and Craig (1994). Rao (1994) was the first to introduce the concept of  $A^*$ -Algebra  $(A, \wedge, \vee, *, (-)\tilde{\cdot}, (-)_{\pi}, 0, 1, 2)$  and also studied its equivalence with ADA, C-algebra, ADA's connection with 3-Ring, stone type representation but also introduced the concept of  $A^*$ -clone, the If-Then-Else structure over  $A^*$ -algebra and Ideal of  $A^*$ -algebra. Venkateswara (2000) introduced the concept of pre  $A^*$ -algebra  $(A, \wedge, \vee, (-)\tilde{\cdot})$  analogous to C-algebra as a reduct of  $A^*$ - algebra. Venkateswara and Srinivasa (2009) defined a partial ordering on a pre  $A^*$ -algebra  $A$  and the

properties of  $A$  as a poset are studied.

## PRELIMINARIES

### Definition 1

Boolean algebra is algebra  $(B, \vee, \wedge, (-)', 0, 1)$  with two binary operations, one unary operation (called complementation), and two nullary operations which satisfies:

- (i)  $(B, \vee, \wedge)$  is a distributive lattice
- (ii)  $x \wedge 0 = 0, x \vee 1 = 1$
- (iii)  $x \wedge x' = 0, x \vee x' = 1$

We can prove that  $x'' = x, (x \vee y)' = x' \wedge y', (x \wedge y)' = x' \vee y'$  for all  $x, y \in B$

An algebra  $(A, \wedge, \vee, (-)\tilde{\cdot})$  satisfying

- (a)  $x\tilde{\cdot}\tilde{\cdot} = x, \forall x \in A,$
- (b)  $x \wedge x = x, \forall x \in A,$
- (c)  $x \wedge y = y \wedge x, \forall x, y \in A,$
- (d)  $(x \wedge y)\tilde{\cdot} = x\tilde{\cdot} \vee y\tilde{\cdot}, \forall x, y \in A,$
- (e)  $x \wedge (y \wedge z) = (x \wedge y) \wedge z, \forall x, y, z \in A$
- (f)  $x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z), \forall x, y, z \in A,$
- (g)  $x \wedge y = x \wedge (x\tilde{\cdot} \vee y), \forall x, y, z \in A,$

is called a Pre  $A^*$ -algebra.

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**Example 1**

$\mathbf{3} = \{0, 1, 2\}$  with operations  $\wedge, \vee, (-)^\sim$  defined below is a pre  $A^*$ -algebra.

$\wedge$	0	1	2	$\vee$	0	1	2	$x$	$x^\sim$
0	0	0	2	0	0	1	2	0	1
1	0	1	2	1	1	1	2	1	0
2	2	2	2	2	2	2	2	2	2

**Note 1**

The elements 0, 1, 2 in the above example satisfy the following laws:

- (a)  $2^\sim = 2$ ; (b)  $1 \wedge x = x$  for all  $x \in \mathbf{3}$ ; (c)  $0 \vee x = x$  for all  $x \in \mathbf{3}$ ; (d)  $2 \wedge x = 2 \vee x = 2$  for all  $x \in \mathbf{3}$ .

**Example 2**

$\mathbf{2} = \{0, 1\}$  with operations  $\wedge, \vee, (-)^\sim$  defined below is a Pre  $A^*$ -algebra.

$\wedge$	0	1	$\vee$	0	1	$x$	$x^\sim$
0	0	0	0	0	1	0	1
1	0	1	1	1	1	1	0

**Note 2**

- (i)  $(\mathbf{2}, \vee, \wedge, (-)^\sim)$  is a Boolean algebra. So every Boolean algebra is a Pre  $A^*$  algebra.
- (ii) The identities  $x^{\sim\sim} = x, \forall x \in A$  and  $(x \wedge y)^\sim = x^\sim \vee y^\sim, \forall x, y \in A$  implies that the varieties of pre  $A^*$ -algebras satisfies all the dual statements of  $x \wedge x = x, \forall x \in A$  to  $x \wedge y = x \wedge (x^\sim \vee y), \forall x, y, z \in A$ .

**Note 3**

Let A be a Pre  $A^*$ -algebra then A is Boolean algebra iff  $x \vee (x \wedge y) = x, x \wedge (x \vee y) = x$  (absorption laws holds).

**Lemma 1**

Every pre  $A^*$ -algebra satisfies the following laws (Venkateswara and Srinivasa, 2009).

- (a)  $x \vee (x^\sim \wedge x) = x$
- (b)  $(x \vee x^\sim) \wedge y = (x \wedge y) \vee (x^\sim \wedge y)$
- (c)  $(x \vee x^\sim) \wedge x = x$
- (d)  $(x \vee y) \wedge z = (x \wedge z) \vee (x^\sim \wedge y \wedge z)$

**Definition 1**

Let A be a Pre  $A^*$ -algebra. An element  $x \in A$  is called central element of A if  $x \vee x^\sim = 1$  and the set  $\{x \in$

$x \vee x^\sim = 1\}$  of all central elements of A is called the centre of A and it is denoted by  $B(A)$ . Note that if A is a pre  $A^*$ -algebra with 1, then  $1, 0 \in B(A)$ . If the centre of pre  $A^*$ -algebra coincides with  $\{0, 1\}$  then we say that A has trivial centre.

**Theorem 1**

Let A be a pre  $A^*$ -algebra with 1, then  $B(A)$  is a Boolean algebra with the induced operations  $\wedge, \vee, (-)^\sim$  (Venkateswara and Srinivasa, 2009).

**Lemma 2**

Let A be a Pre  $A^*$ -algebra with 1 (Venkateswara and Srinivasa, 2009),

- (a) If  $y \in B(A)$  then  $x \wedge x^\sim \wedge y = x \wedge x^\sim, \forall x \in A$
- (b)  $x \wedge (x \vee y) = x \vee (x \wedge y) = x$  if and only if  $x, y \in B(A)$

**SEMILATTICE STRUCTURE ON PRE  $A^*$ -ALGEBRA**

**Theorem 2**

Let A be a Pre  $A^*$ -algebra define a binary operation  $\oplus$  on A by  $x \oplus y = x \vee y$  for all  $x, y \in A$  then  $\langle A, \oplus \rangle$  is a semi lattice.

**Proof**

$x \oplus x = x \vee x = x$  for all  $x \in A$ .  
 For  $x, y \in A$  we have  $x \oplus y = x \vee y = y \vee x = y \oplus x$ .  
 $x \oplus (y \oplus z) = x \oplus (y \vee z)$   
 $= x \vee (y \vee z) = (x \vee y) \vee z = (x \oplus y) \oplus z$ , for all  $x, y, z \in A$ .  
 Hence  $\langle A, \oplus \rangle$  is a semi-lattice.

**Definition 3**

Let A be a pre  $A^*$ -algebra define a relation  $\leq_\oplus$  on A by  $x \leq_\oplus y$  iff  $x \oplus y = y$ .

**Lemma 3**

Let A be a pre  $A^*$ -algebra then  $(A, \leq_\oplus)$  is a poset.

**Proof**

Since  $x \oplus x = x \vee x = x, x \leq_\oplus x$ , for all  $x \in A$   
 Therefore  $\leq_\oplus$  is reflexive.

Suppose  $x \leq_{\oplus} y, y \leq_{\oplus} z$ , for all  $x, y, z \in A$  then  $x \oplus y = y$ , and  $y \oplus z = z$ . Now  $x \oplus z = x \oplus (y \oplus z) = (x \oplus y) \oplus z = y \oplus z = z$  that is  $x \leq_{\oplus} z$ , this shows that  $\leq_{\oplus}$  is Transitive.

Let  $x \leq_{\oplus} y$  and  $y \leq_{\oplus} x$  for all  $x, y \in A$  then  $x \oplus y = y$  and  $y \oplus x = x \Rightarrow x = y$ . This shows that  $\leq_{\oplus}$  is anti symmetric. Therefore  $(A, \leq_{\oplus})$  is a poset.

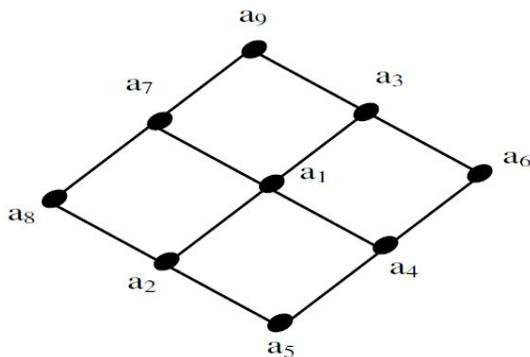
**Note 4**

We have  $x \leq_{\oplus} y$  iff  $x \oplus y = y$ , so  $x \leq_{\oplus} x \oplus y$  for all  $x \in A$ . This shows that  $x \oplus y$  is the supremum of  $\{x, y\}$ .

Let  $A$  be a Pre  $A^*$ -algebra with  $0, 1, 2$  then  $0 \leq_{\oplus} x(x \vee 0 = x$  for all  $x \in A)$  and  $x \leq_{\oplus} 2(2 \vee x = 2$  for all  $x \in A)$ . This gives that  $2$  is the greatest element and  $0$  is the least element of the poset  $(A, \leq_{\oplus})$ . The Hasse diagram of the poset  $(A, \leq_{\oplus})$  is



We have  $A \times A = \{ a_1 = (1,1), a_2 = (1,0), a_3 = (1,2), a_4 = (0,1), a_5 = (0,0), a_6 = (0,2), a_7 = (2,1), a_8 = (2,0), a_9 = (2,2) \}$  is a pre  $A^*$ -algebra under point wise operation and  $A \times A$  is having four central elements and remaining are non central elements, among that  $a_9 = (2,2)$  is satisfying the property that  $a_9 \sim a_9$ . The Hasse diagram is of the poset  $(A \times A, \leq_{\oplus})$  as shown:



Observe that  $x \leq_{\oplus} a_9 (x \vee a_9 = a_9)$  and  $a_5 \leq_{\oplus} x (x \vee a_5 = x)$  for all  $x \in A \times A$ . This shows that  $a_9$  is the greatest element and  $a_5$  is the least element of  $A \times A$ .

**Lemma 4**

The following conditions hold for any elements  $x$  and  $y$  in a pre  $A^*$ -algebra  $A$

- (i)  $x \leq_{\oplus} x \vee y$
- (ii)  $x \wedge x \sim \leq_{\oplus} x \wedge y$

**Proof**

- (i) Consider  $(x \vee y) \oplus x = (x \vee y) \vee x = x \vee y$ . Therefore,  $x \leq_{\oplus} x \vee y$ .
- (ii) Consider  $(x \wedge y) \oplus (x \wedge x \sim) = (x \wedge y) \vee (x \wedge x \sim) = x \wedge (y \vee x \sim) = x \wedge y$  (by dual of  $x \wedge y = x \wedge (x \sim \vee y)$ ),  $\forall x, y, z \in A$

Therefore  $x \wedge x \sim \leq_{\oplus} x \wedge y$

**Lemma 5**

Let  $A$  be a Pre  $A^*$ -algebra then  $\oplus$  is distributive over  $\vee$  and  $\wedge$  that is (i)  $x \oplus (y \vee z) = (x \oplus y) \vee (x \oplus z)$ . (ii)  $x \oplus (y \wedge z) = (x \oplus y) \wedge (x \oplus z)$

**Proof**

- (i)  $(x \oplus y) \vee (x \oplus z) = (x \vee y) \vee (x \vee z) = x \vee (y \vee z) = x \oplus (y \vee z)$
- (ii)  $(x \oplus y) \wedge (x \oplus z) = (x \vee y) \wedge (x \vee z) = x \vee (y \wedge z) = x \oplus (y \wedge z)$

**Theorem 3**

Let  $A$  be a Pre  $A^*$ -algebra for any  $x \in A$  then the following holds in the semi-lattice  $\langle A, \oplus \rangle$ .

- (i)  $x \vee x \sim$  is the supremum of  $\{x, x \sim\}$
- (ii)  $x \wedge x \sim$  is the infimum of  $\{x, x \sim\}$

**Proof**

- (i)  $x \oplus (x \vee x \sim) = x \vee (x \vee x \sim) = x \vee x \sim$

Therefore,  $x \leq_{\oplus} x \vee x^{\sim}$ .

$$x^{\sim} \oplus (x \vee x^{\sim}) = x^{\sim} \vee (x \vee x^{\sim}) = x \vee x^{\sim}.$$

Therefore,  $x^{\sim} \leq_{\oplus} x \vee x^{\sim}$ .

$x \vee x^{\sim}$  is upper bound of  $\{x, x^{\sim}\}$ .

Let  $k$  be the upper bound of  $\{x, x^{\sim}\}$

$$\Rightarrow x \leq_{\oplus} k \text{ and } x^{\sim} \leq_{\oplus} k \text{ that is } x \oplus k = k \text{ and } x^{\sim} \oplus k = k$$

$$\Rightarrow x \vee k = k \text{ and } x^{\sim} \vee k = k.$$

$$\text{Now } k \oplus (x \vee x^{\sim}) = k \vee (x \vee x^{\sim}) = (k \vee x) \vee x^{\sim} \\ = k \vee x^{\sim} = k$$

$$\therefore x \vee x^{\sim} \leq_{\oplus} k.$$

Therefore,  $x \vee x^{\sim}$  is least upper bound of  $\{x, x^{\sim}\}$ .

$$\text{Sup } \{x, x^{\sim}\} = x \vee x^{\sim}$$

$$(ii) \ x \oplus (x \wedge x^{\sim}) = x \vee (x \wedge x^{\sim}) = x$$

Therefore  $x \wedge x^{\sim} \leq_{\oplus} x$ .

$$x^{\sim} \oplus (x \wedge x^{\sim}) = x^{\sim} \vee (x \wedge x^{\sim}) = x^{\sim}$$

Therefore,  $x \wedge x^{\sim} \leq_{\oplus} x^{\sim}$ .

$x \wedge x^{\sim}$  is lower bound of  $\{x, x^{\sim}\}$ .

Let  $l$  be the lower bound of  $\{x, x^{\sim}\}$

$$\Rightarrow l \leq_{\oplus} x \text{ and } l \leq_{\oplus} x^{\sim} \text{ that is } x \oplus l = x \text{ and } x^{\sim} \oplus l = x^{\sim}$$

$$\Rightarrow l \vee x = x \text{ and } l \vee x^{\sim} = x^{\sim}$$

$$\text{Now } l \oplus (x \wedge x^{\sim}) = l \vee (x \wedge x^{\sim}) = (l \vee x) \wedge (l \vee x^{\sim}) \\ = x \wedge x^{\sim} \Rightarrow l \leq_{\oplus} x \wedge x^{\sim}$$

$\therefore x \wedge x^{\sim}$  is greatest lower bound of  $\{x, x^{\sim}\}$ .

$$\text{Inf } \{x, x^{\sim}\} = x \wedge x^{\sim}.$$

### Lemma 6

In the poset  $(A, \leq_{\oplus})$  and  $x, y \in A$ . If  $x \leq_{\oplus} y$  then for  $a \in A$ .

$$(i) \ a \wedge x \leq_{\oplus} a \wedge y$$

$$(ii) \ a \vee x \leq_{\oplus} a \vee y$$

### Proof

If  $x \leq_{\oplus} y$  then  $x \oplus y = y \Rightarrow x \vee y = y$

$$(i) \ (a \wedge x) \oplus (a \wedge y) = (a \wedge x) \vee (a \wedge y) = a \wedge (x \vee y) = a \wedge y$$

$$\therefore a \wedge x \leq_{\oplus} a \wedge y$$

$$(ii) \ (a \vee x) \oplus (a \vee y) = (a \vee x) \vee (a \vee y) = a \vee (x \vee y) = a \vee y$$

$$\therefore a \vee x \leq_{\oplus} a \vee y$$

### Lemma 7

Let  $A$  be a Pre  $A^*$ -algebra for any  $x, y \in A$ ,  $x \leq_{\oplus} x \vee y$  then,  $x \vee y$  is the upper bound of  $\{x, y\}$ .

### Proof

Suppose  $x \leq_{\oplus} x \vee y$  then  $x \vee y$  is upper bound of  $x$

Now  $(x \vee y) \oplus y = (x \vee y) \vee y = x \vee y \Rightarrow y \leq_{\oplus} x \vee y$

Therefore  $x \vee y$  is upper bound of  $y$ .

$\therefore x \vee y$  is the upper bound of  $\{x, y\}$ .

### Theorem 4

Let  $A$  be a Pre  $A^*$ -algebra for any  $x, y \in A$  then  $\text{sup } \{x, y\} = x \vee y$  in the semilattice  $\langle A, \oplus \rangle$ .

### Proof

$$(x \vee y) \oplus x = (x \vee y) \vee x = x \vee y$$

$$\therefore x \leq_{\oplus} x \vee y.$$

$$(x \vee y) \oplus y = (x \vee y) \vee y = x \vee y$$

$$\Rightarrow y \leq_{\oplus} x \vee y$$

Therefore  $x \vee y$  is upper bound of  $y$

$\therefore x \vee y$  is the upper bound of  $\{x, y\}$

Suppose  $m$  is the upper bound of  $\{x, y\}$

$$\Rightarrow x \leq_{\oplus} m \text{ and } y \leq_{\oplus} m \text{ that is } m \oplus x = m \text{ and } m \oplus y = m$$

$$\Rightarrow m \vee x = m \text{ and } m \vee y = m$$

$$\text{Now } m \oplus (x \vee y) = m \vee (x \vee y) = (m \vee x) \vee y = m \vee y = m$$

$$\Rightarrow x \vee y \leq_{\oplus} m$$

$\therefore x \vee y$  is the least upper bound of  $\{x, y\}$

$$\therefore \text{sup } \{x, y\} = x \vee y$$

### Note 5

In general for a pre  $A^*$ -algebra with 1,  $x \vee y$  need not be the greatest lower bound of  $\{x, y\}$  in  $(A, \leq_{\oplus})$ . For example  $2 \vee x = 2 \wedge x = 2, \forall x \in A$  is not a greatest lower bound. However we have the following theorem.

### Theorem 5

In a semi lattice  $\langle A, \oplus \rangle$  with 1, for any  $x, y \in B(A)$  then  $\text{inf } \{x, y\} = x \wedge y$

### Proof

If  $x, y \in B(A)$ , then by lemma 2b,  $x \vee (x \wedge y) = x$  and  $y \vee (x \wedge y) = y$

$$\Rightarrow x \oplus (x \wedge y) = x \text{ and } y \oplus (x \wedge y) = y$$

This shows that  $x \wedge y \leq_{\oplus} x$  and  $x \wedge y \leq_{\oplus} y$

Hence  $x \wedge y$  is a lower bound of  $\{x, y\}$ .

Suppose  $k$  is a lower bound of  $\{x, y\}$ , then  $k \leq_{\oplus} x, k \leq_{\oplus} y$

$$\Rightarrow k \oplus x = x \text{ and } k \oplus y = y$$

$$\Rightarrow x \vee k = x, y \vee k = y$$

Now  $k \oplus (x \wedge y) = k \vee (x \wedge y) = (k \vee x) \wedge (k \vee y) = x \wedge y$ .

Therefore,  $k \leq_{\oplus} x \wedge y$

$x \wedge y$  is the greatest lower bound of  $\{x, y\}$ . Hence,  $\inf \{x, y\} = x \wedge y$

### Theorem 6

If  $A$  is a Pre  $A^*$ -algebra and  $x \vee (x \wedge y) = x$ , for all  $x, y \in A$  then  $(A, \leq_{\oplus})$  is a lattice.

### Proof

By theorem 4 every pair of elements have supremum. If  $x \vee (x \wedge y) = x$  for all  $x, y \in A$  then by theorem 5 every pair of elements have infimum. Hence  $(A, \leq_{\oplus})$  is a lattice.

### Lemma 8

Let  $A$  be a pre  $A^*$ -algebra then:

- (i)  $x \vee (x \oplus y) = x \vee y$ .
- (ii)  $(x \oplus y) \vee x = x \oplus y$ .

### Proof

- (i)  $x \vee (x \oplus y) = x \vee (x \vee y) = x \vee y$
- (ii)  $(x \oplus y) \vee x = (x \vee y) \vee x = x \vee y = x \oplus y$

Now we present a number of equivalent conditions for a pre  $A^*$ -algebra become a Boolean algebra.

### Theorem 7

The following conditions are equivalent for any pre  $A^*$ -algebra  $(A, \wedge, \vee, (-)^{\sim})$

- (1)  $A$  is Boolean algebra
- (2)  $x \wedge y \leq_{\oplus} x$  for all  $x, y \in A$
- (3)  $x \wedge y \leq_{\oplus} y$  for all  $x, y \in A$
- (4)  $x \wedge y$  is a lower bound of  $\{x, y\}$  in  $(A, \leq_{\oplus})$  for all  $x, y \in A$
- (5)  $x \wedge y$  is a infimum of  $\{x, y\}$  in  $(A, \leq_{\oplus})$  for all  $x, y \in A$
- (6)  $x \vee x^{\sim}$  is the least element in  $(A, \leq_{\oplus})$  for every  $x \in A$

### Proof

(1)  $\Rightarrow$  (2) Suppose  $A$  is a Boolean algebra.

Now  $x \oplus (x \wedge y) = x \vee (x \wedge y) = x$  (by absorption law)

$\therefore x \wedge y \leq_{\oplus} x$

(2)  $\Rightarrow$  (3) Suppose  $x \wedge y \leq_{\oplus} x$  then  $x \oplus (x \wedge y) = x$ .

Therefore  $x \vee (x \wedge y) = x$

Now  $y \oplus (x \wedge y) = y \vee (x \wedge y) = y$ . Therefore,  $x \wedge y \leq_{\oplus} y$

(3)  $\Rightarrow$  (4) Suppose that  $x \wedge y \leq_{\oplus} y \Rightarrow y \oplus (x \wedge y) = y$

therefore  $y \vee (x \wedge y) = y$

Since  $x \wedge y \leq_{\oplus} y$  then  $x \wedge y$  is lower bound of  $y$

Now  $x \oplus (x \wedge y) = x \vee (x \wedge y) = x$  (by supposition)

$\therefore x \wedge y \leq_{\oplus} x$

$\Rightarrow x \wedge y$  is a lower bound of  $x$ .

$\therefore x \wedge y$  is a lower bound of  $\{x, y\}$ .

(4)  $\Rightarrow$  (5) Suppose  $x \wedge y$  is a lower bound of  $\{x, y\}$

Suppose  $z$  is a lower bound of  $\{x, y\}$  then  $z \leq_{\oplus} x, z \leq_{\oplus} y$  that is  $x \oplus z = x$  and  $y \oplus z = y$

$\Rightarrow x \vee z = x, y \vee z = y$

Now  $z \oplus (x \wedge y) = z \vee (x \wedge y) = (z \vee x) \wedge (z \vee y)$

$= x \wedge y$

Therefore,  $z \leq_{\oplus} x \wedge y$

$x \wedge y$  is the greatest lower bound of  $\{x, y\}$

Hence  $\inf \{x, y\} = x \wedge y$

(5)  $\Rightarrow$  (6) Suppose  $\inf \{x, y\} = x \wedge y$  then  $x, y \in B(A)$

Now  $\inf \{x \wedge x^{\sim}, y\} = x \wedge x^{\sim} \wedge y = x \wedge x^{\sim}$  (by lemma 2a)

$\Rightarrow x \wedge x^{\sim} \leq_{\oplus} y$ . Therefore  $x \wedge x^{\sim}$  is the least element in  $(A, \leq_{\oplus})$ .

(6)  $\Rightarrow$  (1) Suppose  $x \wedge x^{\sim}$  is the least element in  $A$  then

$x \wedge x^{\sim} \leq_{\oplus} y$ , for  $y \in A$

$\Rightarrow (x \wedge x^{\sim}) \oplus y = y \Rightarrow (x \wedge x^{\sim}) \vee y = y$

Now  $y \wedge (x \vee y) = [(x \wedge x^{\sim}) \vee y] \vee (x \vee y)$

$= [(x \wedge x^{\sim}) \vee x] \vee y = (x \wedge x^{\sim}) \vee y = y$  (by supposition).

Therefore by Note 3 we have  $B$  is Boolean algebra.

### Theorem 8

Let  $A$  be a pre  $A^*$ -algebra  $x \vee x^{\sim}$  is the greatest element in  $(A, \leq_{\oplus})$  for every  $x \in A$  then  $A$  is Boolean algebra.

### Proof

Suppose  $x \vee x^{\sim}$  is the greatest element in  $(A, \leq_{\oplus})$  then  $y$

$\leq_{\oplus} x \vee x^{\sim}$

$\Rightarrow (x \vee x^{\sim}) \oplus y = x \vee x^{\sim}$

$\Rightarrow (x \vee x^{\sim}) \vee y = x \vee x^{\sim}$

Now  $x \vee (x \wedge y) = [x \wedge (x^{\sim} \vee x)] \vee (x \wedge y)$

$= x \wedge [(x \vee x^{\sim}) \vee y] = x \wedge (x \vee x^{\sim}) = x$  (by supposition)

$\therefore x \vee (x \wedge y) = x$ , absorption law holds

By Note 3 we have  $B$  is Boolean algebra.

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